

An Interval Method for Linear IVPs for ODEs

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Improved version of the talk given at the
Workshop on Taylor Models
17–20 December 2003, Miami, Florida

The Problem

Enclose the solution of a system of $n \geq 2$ equations IVP

$$y' = A(t)y + g(t), \quad y(0) = y_0 \in [y_0].$$

Idea (Lohner, Nickel)

- Perform $(n + 1)$ integrations of points specifying a parallelepiped at t_i and enclose each point solution at t_{i+1} .
We have $(n + 1)$ boxes.
- Find $(n + 1)$ points that determine a parallelepiped, which encloses all the parallelepipeds with vertices in these boxes.
- Repeat.

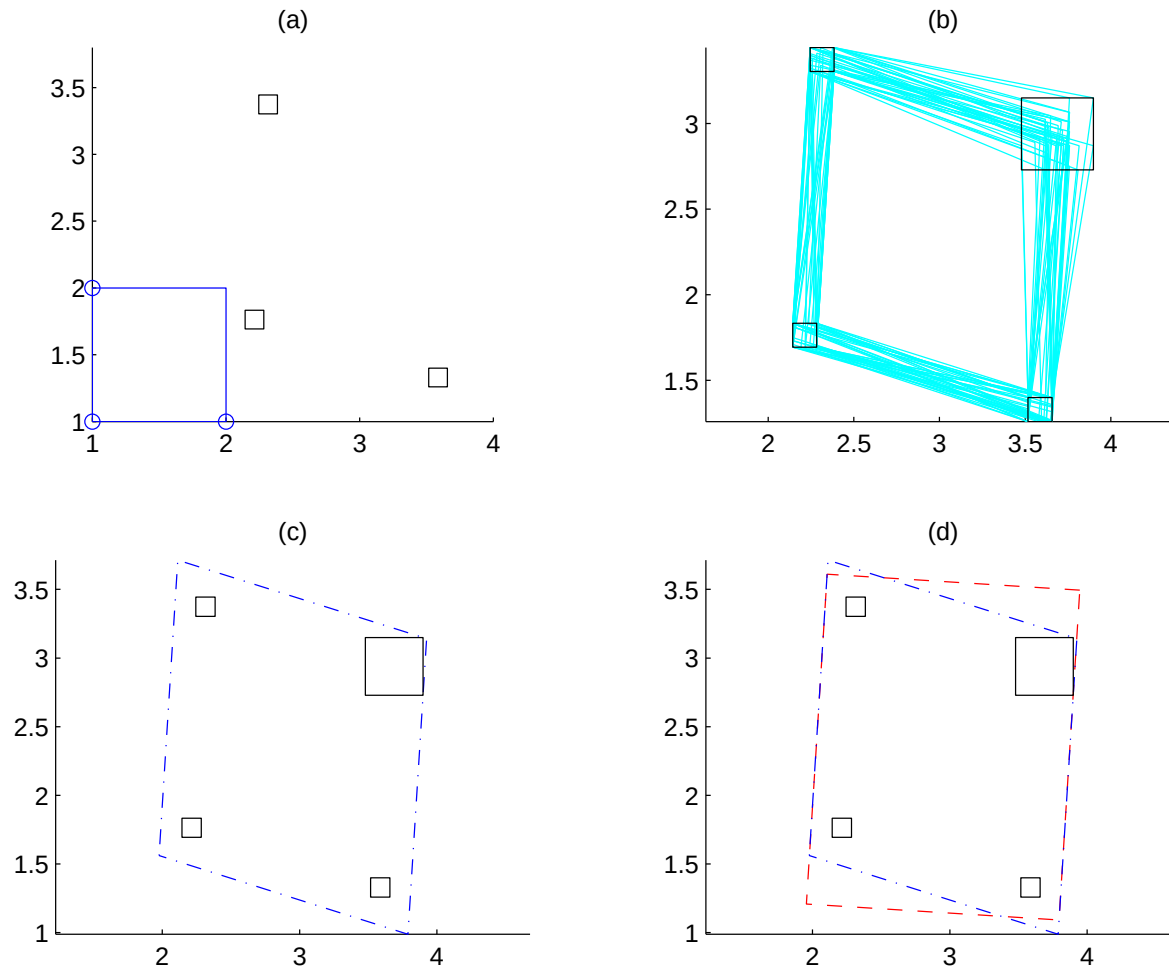


Figure 1: (a) enclosures of point solutions at t_1 ; (b) some of the parallelepipeds with vertices in these enclosures (boxes); the larger box contains the fourth vertices; (c–d) parallelepipeds enclosing the true solution

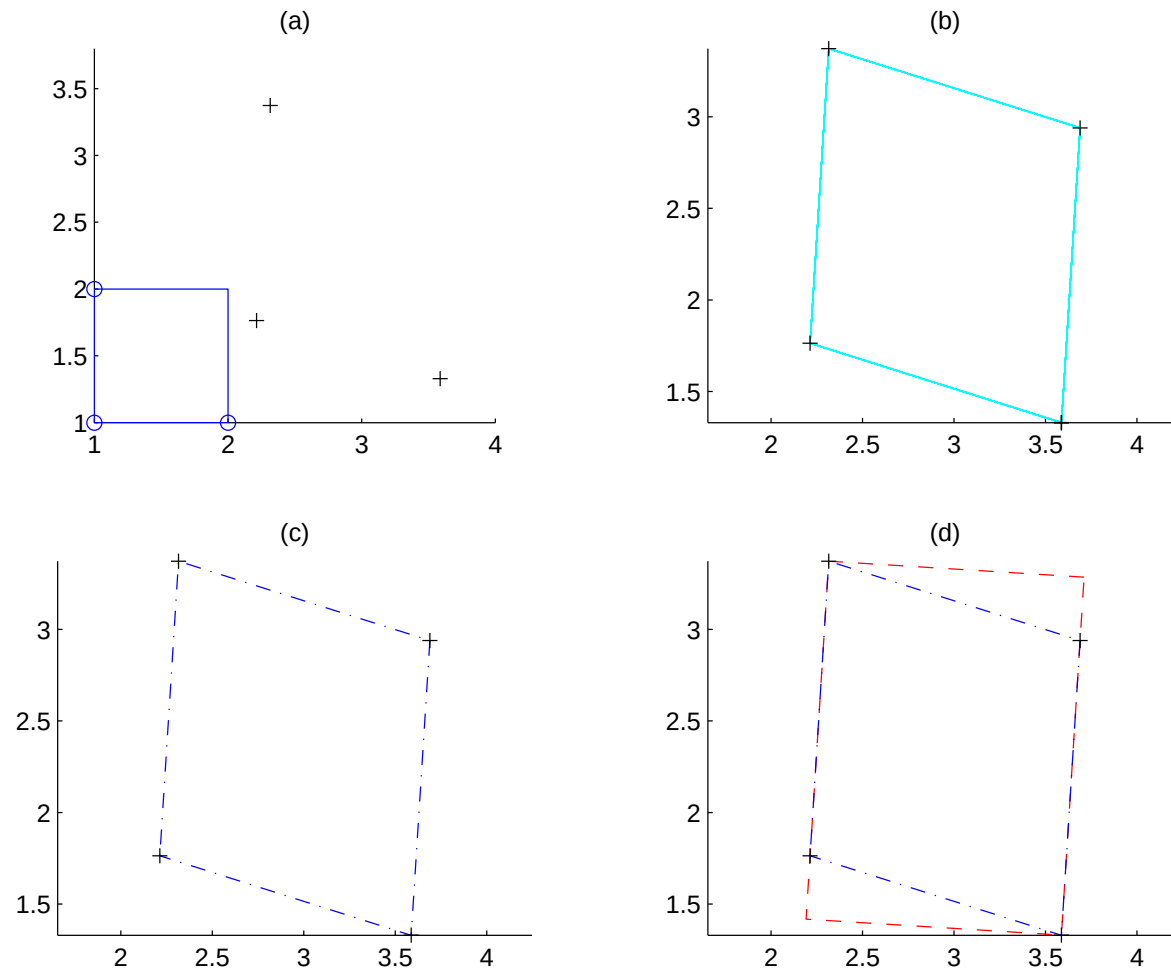


Figure 2: The same computation as in the previous figure, except that the width of each component of the enclosures is 2×10^{-10} . The boxes are denoted by “+”.

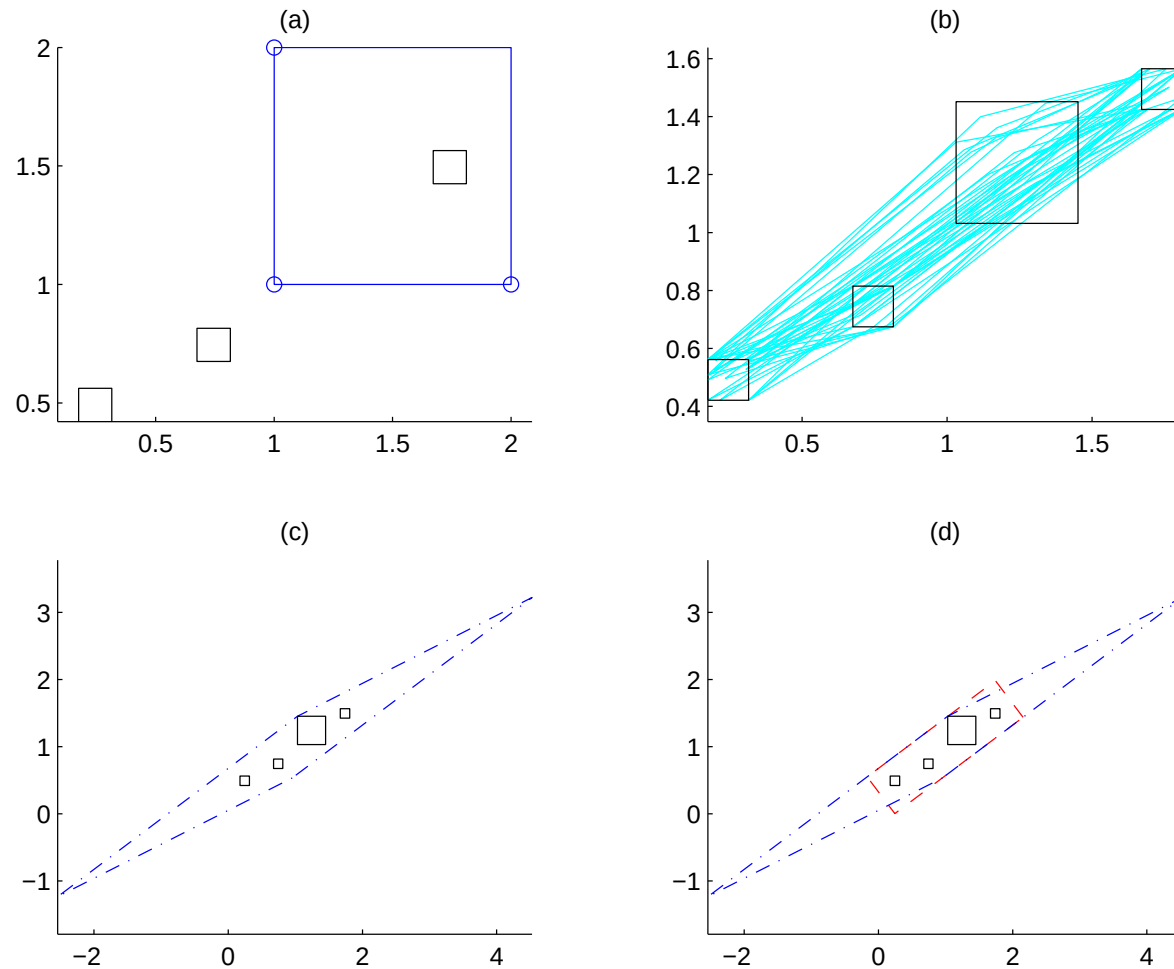


Figure 3: (a) enclosures of point solutions at t_1 ; (b) some of the parallelepipeds with vertices in these enclosures (boxes); the larger box contains the fourth vertices; (c–d) parallelepipeds enclosing the true solution

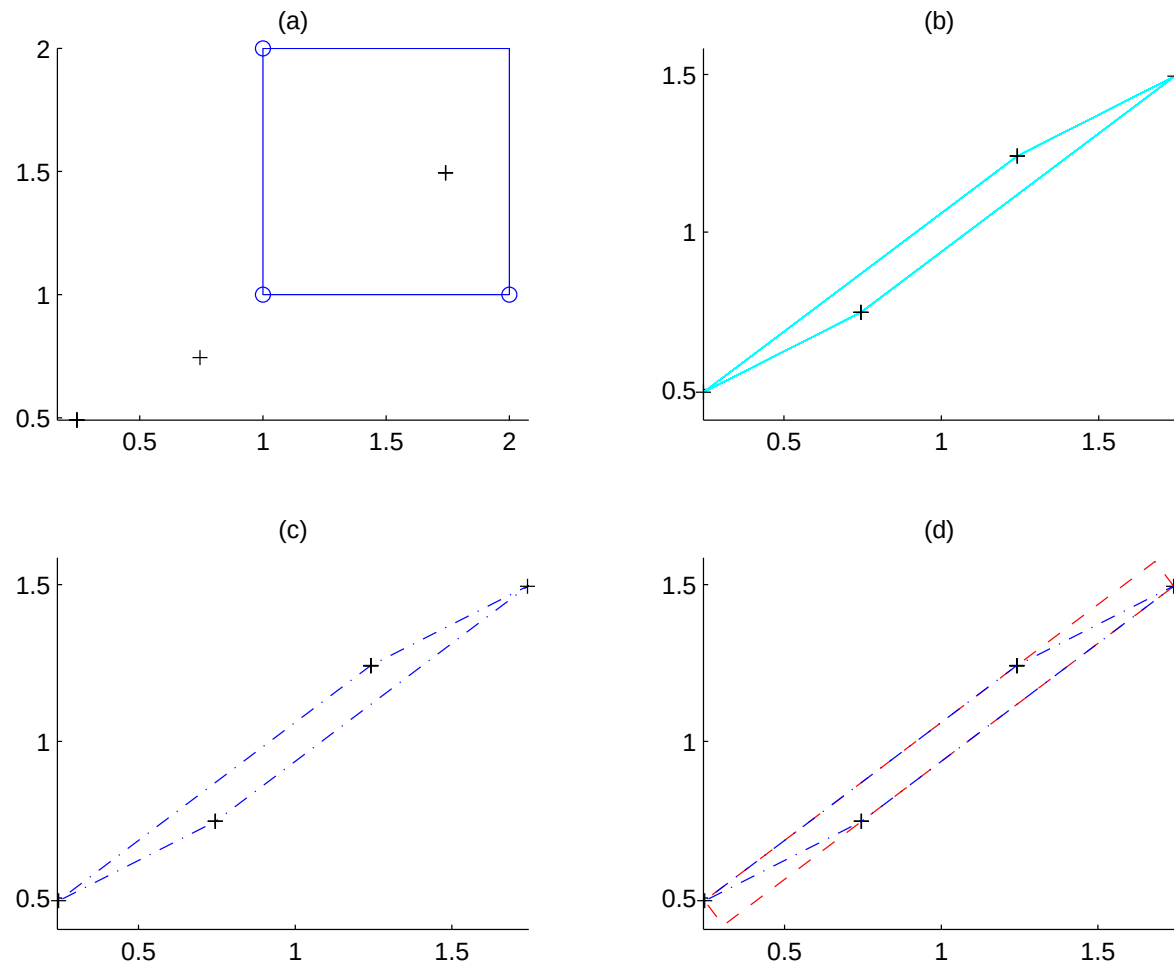


Figure 4: The same computation as in the previous figure, except that the width of each component of the enclosures is 2×10^{-10} . The boxes are denoted by “+”.

Advantages

- We enclose point solutions:
Taylor series + remainder term.
- The method does not impose restrictions on the size of the initial box.
- An automatic differentiation package for computing Taylor coefficients for the solution to $Y' = A(t)Y$, $Y(0) = I$ is not needed.
These coefficients are computed in AWA and VNODE.

Difficulties

- How to compute $(n + 1)$ points on each step such that the parallelepiped specified by them encloses the solution set.
- How to achieve small overestimations and reduce the wrapping effect.

Outline

1. Enclosing point solutions
2. Computing a parallelepiped
3. Choice of a transformation matrix
4. Reducing the wrapping effect
5. Concluding remarks

Enclosing Point Solutions

Denote by $f_i(\cdot)$ the i th Taylor coefficient of the solution to

$$y' = A(t)y + g(t). \quad (1)$$

If h and $[\tilde{y}_0] \ni y_0$ are such that

$$y_0 + \sum_{i=1}^{p-1} t^i f_i(y_0) + t^p f_p([\tilde{y}_0]) \subseteq [\tilde{y}_0] \quad \text{for all } t \in [0, h],$$

then (1) with $y(0) = y_0$ has a unique solution in $[0, h]$, and

$$y(t; t_0, y_0) \in [\tilde{y}_0] \quad \text{for all } t \in [0, h].$$

At $t = h$,

$$y(h; t_0, y_0) \in y_0 + \sum_{i=1}^{p-1} h^i f_i(y_0) + h^p f_p([\tilde{y}_0]).$$

Assume that at a point t_i , for all $y_0 \in [y_0]$,

$$y(t_i; t_0, y_0) \in \{ b_0 + B\alpha \mid \alpha \in [0, 1]^n \},$$

where $B \in \mathbb{R}^{n \times n}$, and $[0, 1]^n$ denotes the vector with each component $[0, 1]$.

We integrate $v_0 = b_0, v_1 = b_0 + b_1, \dots, v_n = b_0 + b_n$
to compute $[w_0], [w_1], \dots, [w_n]$.

That is, for each v_j ,

$$y(t_{i+1}; t_i, v_j) \in [w_j] = v_j + \sum_{i=1}^{p-1} h^i f_i(v_j) + h^p f_p([\tilde{v}_j]),$$

where $v_j \in [\tilde{v}_j]$.

Computing a Parallelepiped

Denote

$$c_j = \text{mid}([w_j]),$$

$$[e_j] = [w_j] - c_j, \quad j = 0, \dots, n,$$

C the $n \times n$ matrix with j th column $c_j - c_0$, and

$$[e] = \sum_{j=1}^n [e_j] + (n-1)[e_0].$$

For all $y_i \in \{b_0 + B\alpha \mid \alpha \in [0, 1]^n\}$,

$$\begin{aligned} y(t_{i+1}; t_i, y_i) &\in \left\{ w_0 + \sum_{j=1}^n \alpha_j (w_j - w_0) \mid \alpha_j \in [0, 1], w_j \in [w_j] \right\} \\ &\subseteq \{ c_0 + C\alpha + [e] \mid \alpha \in [0, 1]^n \}, \end{aligned}$$

since for $\alpha \in [0, 1]^n$, $w_j \in [w_j]$, and $e_j = w_j - c_j \in [e_j]$ ($j = 0, \dots, n$),

$$\begin{aligned}
& w_0 + \sum_{j=1}^n \alpha_j (w_j - w_0) \\
&= c_0 + C\alpha + (w_0 - c_0) + \sum_{j=1}^n \alpha_j (w_j - w_0 - (c_j - c_0)) \\
&= c_0 + C\alpha + e_0 + \sum_{j=1}^n \alpha_j (e_j - e_0) \\
&= c_0 + C\alpha + \sum_{j=1}^n \alpha_j e_j + (1 - \sum_{j=1}^n \alpha_j) e_0 \\
&\in \left\{ c_0 + C\alpha + \sum_{j=1}^n [e_j] + (n-1)[e_0] \mid \alpha \in [0, 1]^n \right\} \\
&= \left\{ c_0 + C\alpha + [e] \mid \alpha \in [0, 1]^n \right\}.
\end{aligned}$$

(Note that each $[e_j]$ is symmetric.)

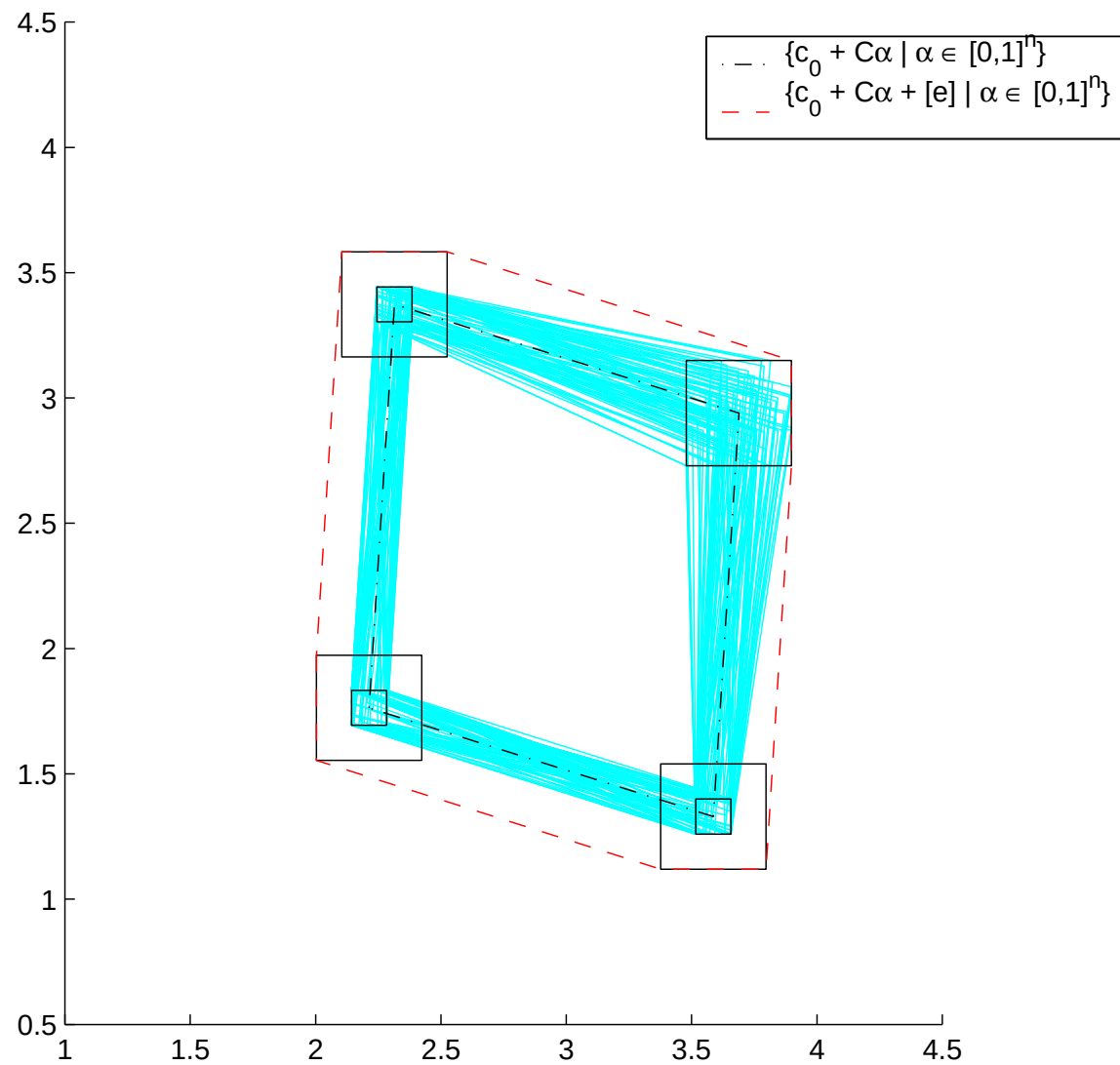


Figure 5: We want to enclose the set $\{c_0 + C\alpha + [e] \mid \alpha \in [0,1]^n\}$ by a parallelepiped.

We want to find g_0 and G such that

$$\{c_0 + C\alpha + e \mid \alpha \in [0, 1]^n, e \in [e]\} \subseteq \{g_0 + G\alpha \mid \alpha \in [0, 1]^n\}.$$

Let $H \in \mathbb{R}^{n \times n}$ be nonsingular.

Denote

$$[r] = (H^{-1}C)[0, 1]^n + H^{-1}[e] \quad \text{and} \quad D = \text{diag}(w([r])).$$

Then

$$\begin{aligned} & \{c_0 + C\alpha + e \mid \alpha \in [0, 1]^n, e \in [e]\} \\ &= \{c_0 + H((H^{-1}C)\alpha + H^{-1}e) \mid \alpha \in [0, 1]^n, e \in [e]\} \\ &\subseteq \{c_0 + Hr \mid r \in [r]\} \\ &= \{c_0 + H\underline{r} + Hr \mid r \in [0, \bar{r} - \underline{r}] = D[0, 1]^n\} \\ &= \{(c_0 + H\underline{r}) + (HD)\alpha \mid \alpha \in [0, 1]^n\} \\ &= \{g_0 + G\alpha \mid \alpha \in [0, 1]^n\}. \end{aligned}$$

This derivation is by R. Lohner (2001, private communications).

Now, for all $y_i \in \{ b_0 + B\alpha \mid \alpha \in [0, 1]^n \}$,

$$y(t_{i+1}; t_i, y_i) \in \{ g_0 + G\alpha \mid \alpha \in [0, 1]^n \}.$$

We integrate $g_0, (g_0 + g_1), \dots, (g_0 + g_n)$.

Subtlety: we compute in floating-point arithmetic $\tilde{g}_0$ and $\tilde{G}$ corresponding to g_0 and G .

Is

$$\{ c_0 + C\alpha + e \mid \alpha \in [0, 1]^n, e \in [e] \} \subseteq \{ \tilde{g}_0 + \tilde{G}\alpha \mid \alpha \in [0, 1]^n \} ? \quad (2)$$

If

$$\tilde{G}^{-1}(c_0 - g_0) + (\tilde{G}^{-1}C)[0, 1]^n + \tilde{G}^{-1}[e] \subseteq [0, 1]^n \quad (3)$$

then (2) holds.

If (3) does not hold in computer arithmetic, inflate $[e]$ and try again.

Choice of a Transformation Matrix

Parallelepiped method

$$H = C,$$

$$[r] = (H^{-1}C)[0, 1]^n + H^{-1}[e] = [0, 1]^n + C^{-1}[e].$$

This method breaks down when C is close to singular.

QR-factorization method

$$C = QR, \quad H = Q,$$

$$[r] = (H^{-1}C)[0, 1]^n + H^{-1}[e] = R[0, 1]^n + Q^T[e].$$

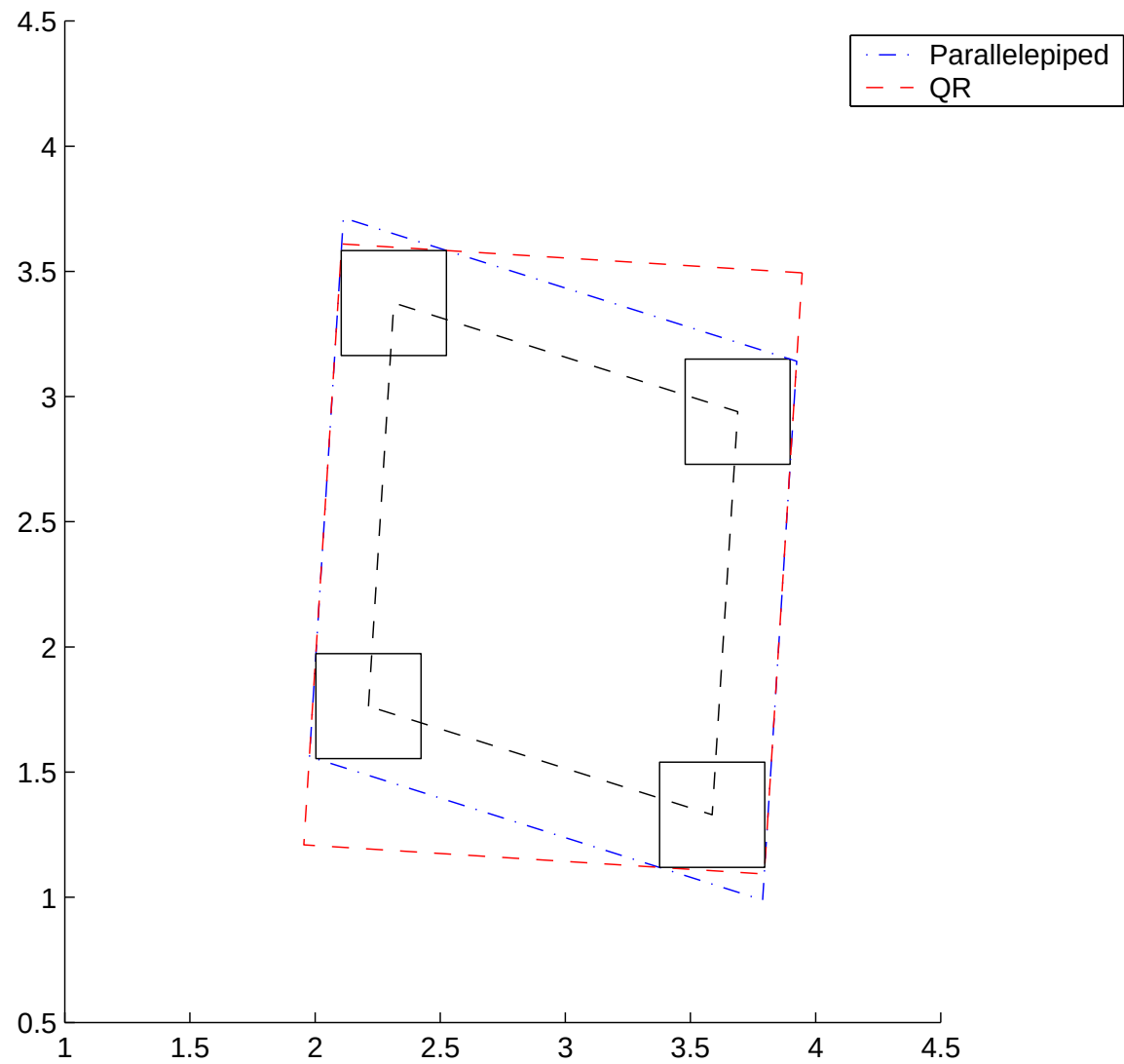


Figure 6: Enclosures obtained by the parallelepiped and QR approaches.

On some problems, with a large initial box, the QR method can produce large overestimations.

Example:

$$y' = \begin{pmatrix} 1 & -2 \\ 3 & -4 \end{pmatrix} y, \quad y(0) \in ([1, 2], [1, 2])^T.$$

We take $[e] = [-10^{-3}, 10^{-3}]$, $h = 0.2$.

The eigenvalues of $\exp(hA)$ are ≈ 0.8187 and 0.6703 .

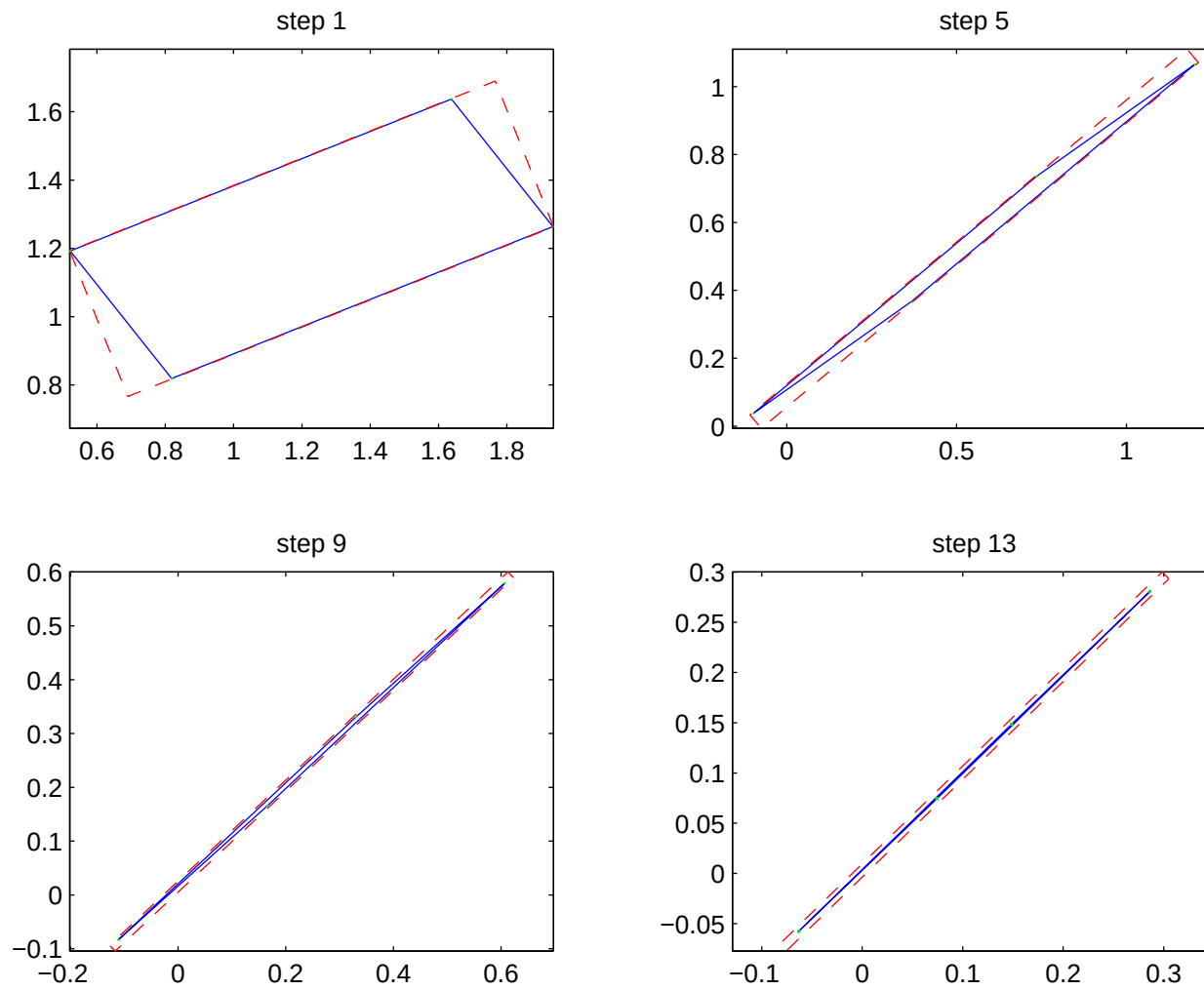


Figure 7: QR; the blue lines denote the true solution set.

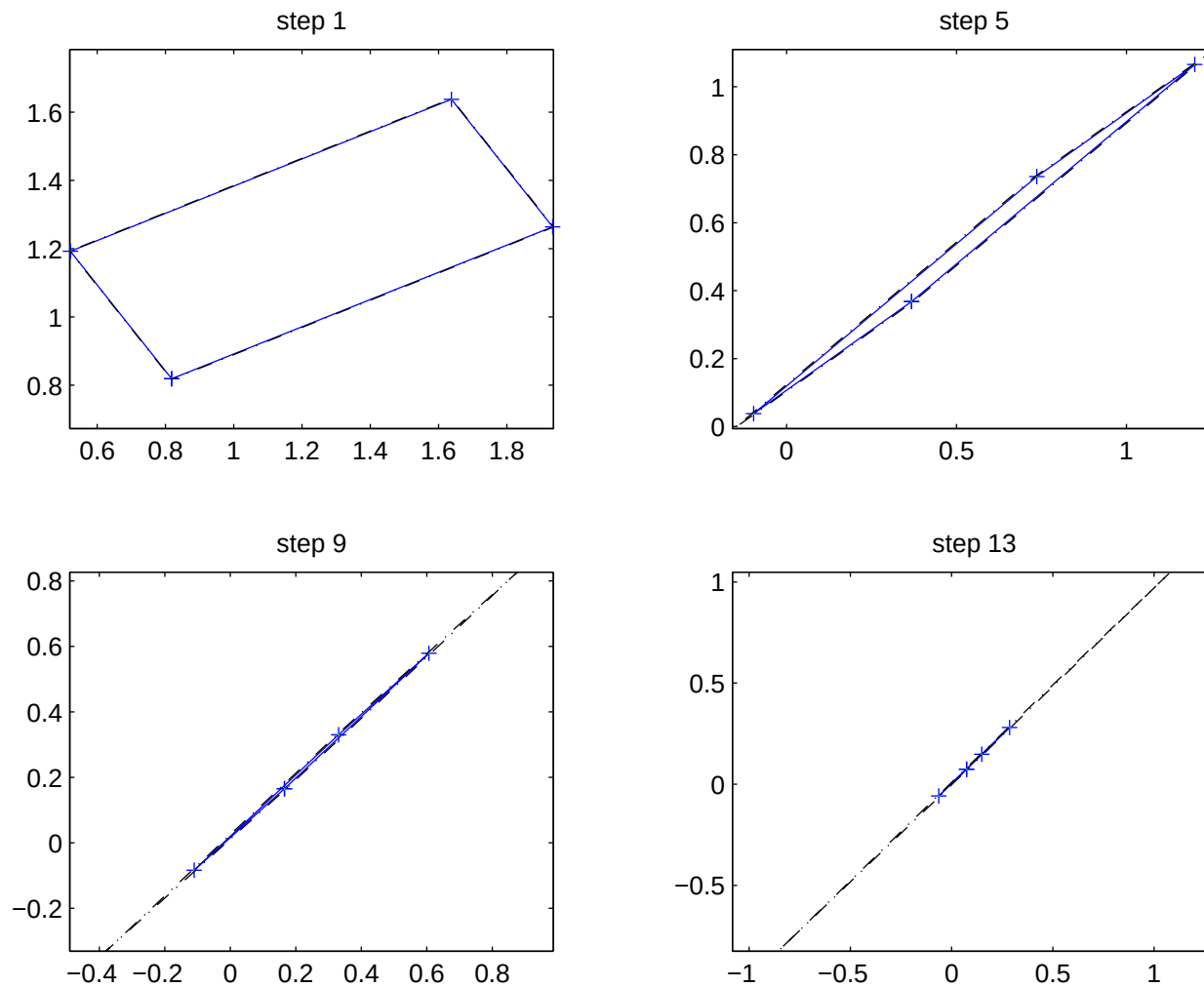


Figure 8: Parallelepiped; the vertices of the true solution are denoted by “+”.

Example:

$$y' = \begin{pmatrix} -0.5 & 1 \\ -1 & 0 \end{pmatrix} y, \quad y(0) \in ([1, 2], [1, 2])^T.$$

We take $[e] = [-10^{-4}, 10^{-4}]$, $h = 0.2$.

The eigenvalues of $\exp(hA)$ are $\approx 0.9334 \pm 0.1831i$, and $\rho(\exp(hA)) \approx 0.9512$.

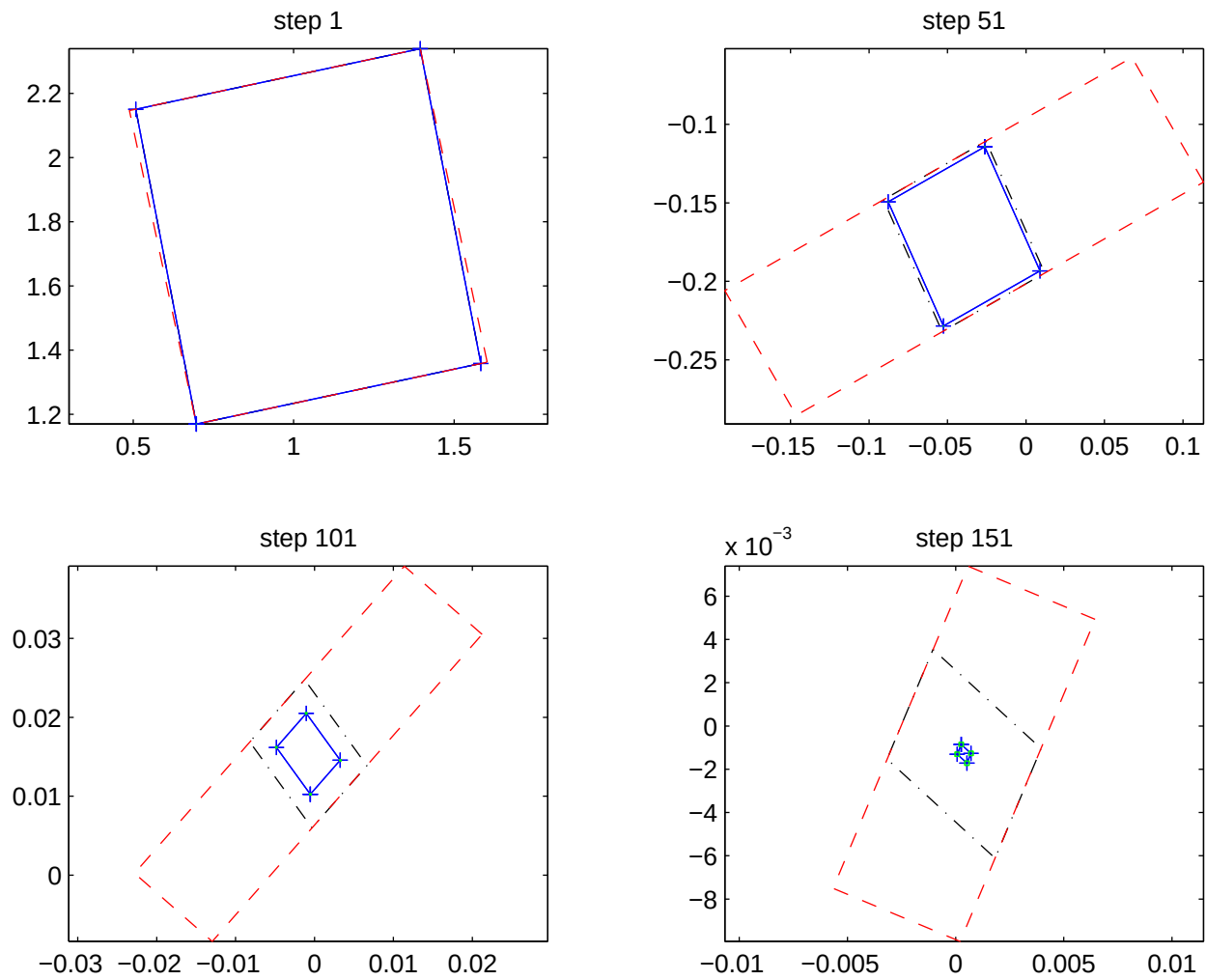


Figure 9: QR and parallelepiped methods

Reducing the Wrapping Effect

The true solution is in

$$\begin{aligned} & \{ g_0 + G\alpha \mid \alpha \in [0, 1]^n \} \\ &= \{ c_0 + Hr \mid r \in (H^{-1}C)[0, 1]^n + H^{-1}[e] \}. \end{aligned}$$

Parallelepiped

$$H = C, \quad [r_P] = [0, 1]^n + C^{-1}[e].$$

QR factorization

$$C = QR, \quad H = Q, \quad [r_Q] = R[0, 1]^n + Q^T[e].$$

Can we combine them, or switch between them at run time?

Two ad-hoc solutions: Approach I and II.

Approach I

We can (roughly) measure the overestimations in the parallelepiped and QR methods by $\|w(C[r_P])\|$ and $\|w(Q[r_Q])\|$, respectively.

Select:

if $\|w(C[r_P])\| \leq \|w(Q[r_Q])\|$

$H = C, [r] = [r_P]$ (parallelepiped)

else

$H = Q, [r] = [r_Q]$ (QR)

Example:

$$y' = \begin{pmatrix} 1 & -2 \\ 3 & -4 \end{pmatrix} y, \quad y(0) \in ([1, 2], [1, 2])^T,$$

$$[e] = [-10^{-3}, 10^{-3}], \quad h = 0.2.$$

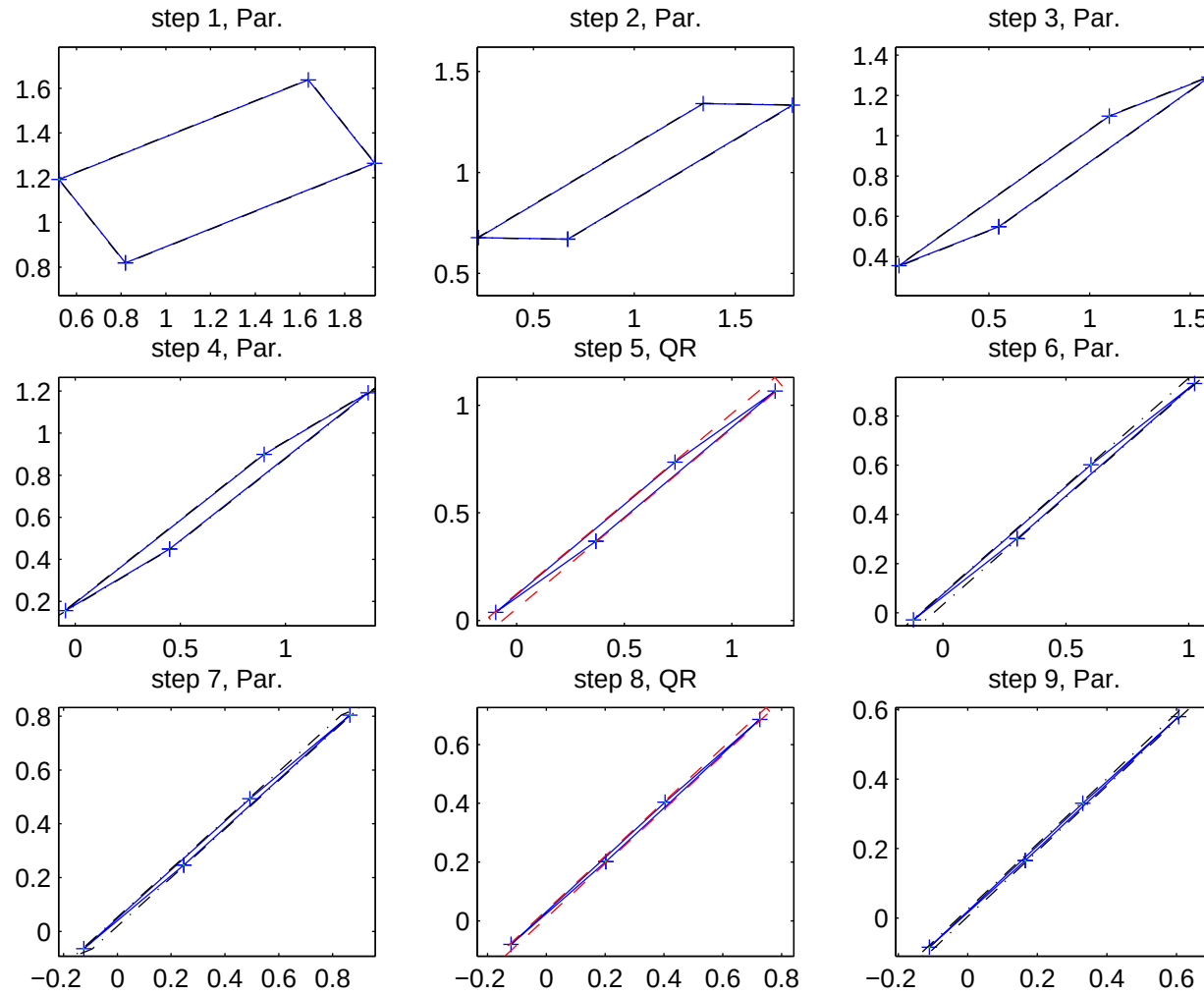


Figure 10: Approach I

Approach II

Let $\beta_{\max}$ be the largest angle among the angles between every two columns of C .

Let $\beta_{\min}$ be the smallest such angle.

Let θ , $0 < \theta \ll \pi$, be a constant.

Select:

if $\beta_{\min} > \theta$ and $\beta_{\max} < \pi - \theta$

$H = C$, $[r] = [r_P]$ (parallelepiped)

else

$H = Q$, $[r] = [r_Q]$ (QR)

Example:

$$y' = \begin{pmatrix} 1 & -2 \\ 3 & -4 \end{pmatrix} y, \quad y(0) \in ([1, 2], [1, 2])^T,$$

$$[e] = [-10^{-3}, 10^{-3}], \quad h = 0.2,$$

$$\theta = 10^\circ = \pi/18.$$

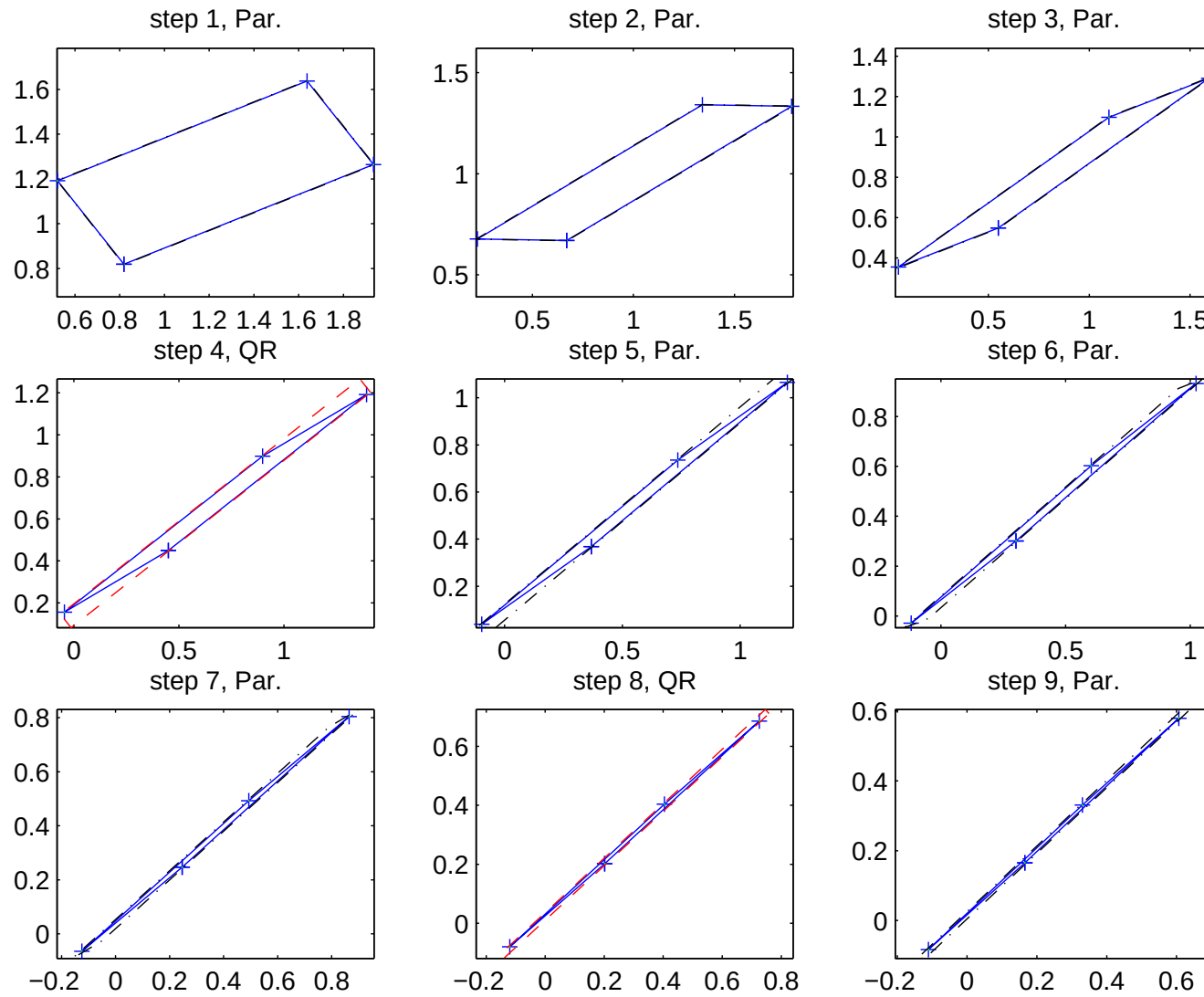


Figure 11: Approach II

Concluding Remarks

- To reduce the wrapping effect when propagating larger sets, a **combination** of the parallelepiped and QR-factorization methods may be necessary.
- When to switch from one method to the other?
- An eigenvalue, or stability type analysis of a combined approach may be necessary.