

A priori Enclosures for ODEs: An Alternative to Fixed Point Iterations

Markus Neher
Institut für Angewandte Mathematik
Universität Karlsruhe

- I. Introduction
- II. Linear ODEs
- III. Numerical Examples
- IV. Nonlinear ODEs

Introduction

Smooth IVP: $u' = f(t, u), \quad u(t_0) = u_0.$

$$(u = (u_1, \dots, u_n), \quad f = (f_1, \dots, f_n)).$$

Moore's enclosure method:

- Automatic computation of Taylor coefficients.
- Interval iteration: For $j := 0, 1, \dots$:

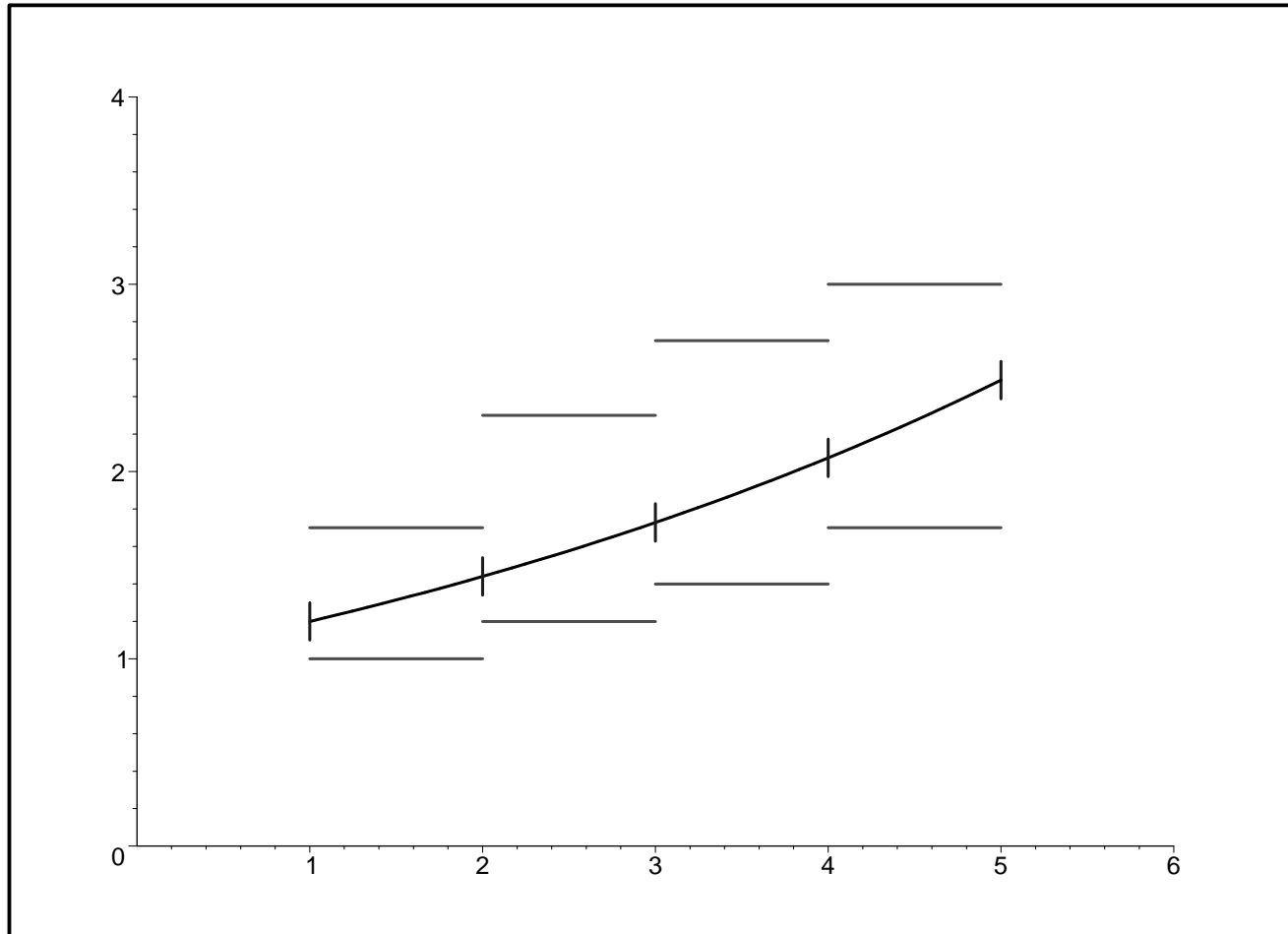
A priori enclosure: $[\hat{u}_{j+1}] \supseteq u(t)$ for all $t \in [t_j, t_{j+1}]$ ("Algorithm I").

Truncation error: $[z_{j+1}] := \frac{h^{(m+1)}}{(m+1)!} f^{(m)}([t_j, t_{j+1}], [\hat{u}_{j+1}]).$

$$u(t_{j+1}) \in [u_{j+1}] := [u_j] + \sum_{k=1}^m \frac{h^k}{k!} f^{(k-1)}(t_j, [u_j]) + [z_{j+1}]$$

("Algorithm II").

Piecewise constant a priori enclosure



A priori Enclosures

- Moore:

Fixed point iteration with constant enclosures:

For some $h > 0$ determine interval $[\hat{u}_1]$ such that

$$u_0 + [0, h] \cdot f([t_0, t_1], [\hat{u}_1]) \subseteq [\hat{u}_1]$$

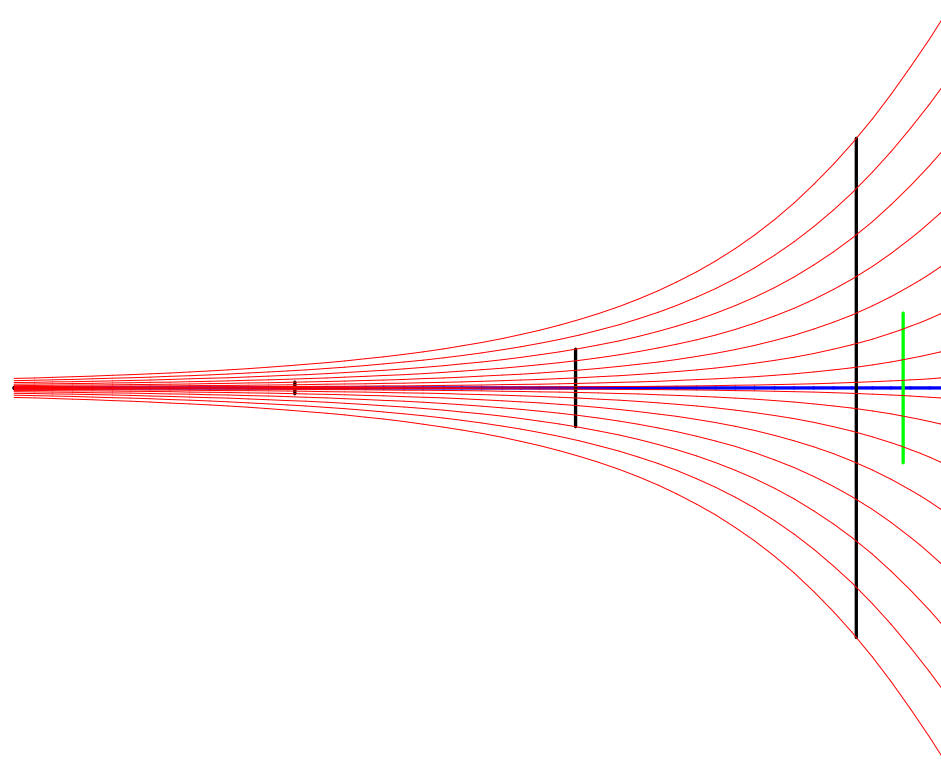
⇒ Step size restrictions: Explicit Euler steps.

- Improvements suggested by Lohner, Corliss & Rihm, Nedialkov & Jackson,
- Berz & Makino: High order TM enclosure by FP iteration

Step Size Control

- Maximum step size for practical calculations:
Slightly less than the radius of convergence of the Taylor series.
- Large step sizes expected to reduce overall wrapping operations.
- Helpful for tracing unstable solutions.

Tracing Unstable Solutions



Linear n -th Order ODE with Analytic Coefficient Functions

$$y^{(n)} = \sum_{i=0}^{n-1} p_i(x) y^{(i)} + p_{-1}(x), \quad x \in [0, r],$$

$$y^{(i)}(0) = y_{i0}, \quad i = 0, 1, \dots, n-1,$$

$$p_i(x) = \sum_{j=0}^{\infty} b_{ij} x^j, \quad |x| < R \quad (R > r), \quad i = -1, 0, \dots, n-1.$$

Power series solution: $y(x) := \sum_{k=0}^{\infty} a_k x^k, \quad |x| < R.$

$$a_{k+n} = \sum_{i=0}^{n-1} \sum_{j=0}^k \frac{P(k-j, i) b_{ij}}{P(k, n)} a_{k+i-j} + \frac{b_{-1, k}}{P(k, n)} \quad \text{for } k = 0, 1, \dots,$$

where $P(k, i) := (k+1) \cdots (k+i), \quad P(k, 0) := 1 \quad \text{for } i \in \mathbb{N}, k \in \mathbb{N}_0.$

Estimation of Recurrence Relation

I. Required: $\forall i \exists m_i \geq i + 2, \exists B_i \geq 0$ so that

$$|b_{ij}| \leq \frac{B_i}{P(j - m_i, m_i)r^j} \quad \text{for all } j > m_i.$$

II. Find $K \geq mn, q \in [0, 1]$:

$$\frac{\max_{\nu=K-m}^{K+n-1} |a_\nu r^\nu|}{\max_{\nu=0}^{K+n-1} |a_\nu r^\nu|} \leq q$$

Estimation of Recurrence Relation

III. Recess condition: Find κ such that

$$\begin{aligned}
 & q \sum_{i=0}^{n-1} \sum_{j=0}^{m_i} \frac{r^{n-i+j} P(\kappa - j, i) |b_{ij}|}{P(\kappa, n)} + \frac{B_{-1} r^{n+m-1}}{P(\kappa - m_{-1}, m_{-1} + n) \max_{\nu=0}^{\kappa+n-1} |a_\nu r^\nu|} \\
 & + \sum_{i=0}^{n-1} \frac{B_i r^{n+m_i-i}}{P(\kappa, n)} \left(\frac{(-1)^{i+1} i!}{P(m_i - 2 - i, i + 1) P(\kappa - m_i + 1 + i, m_i - 1 - i)} \right. \\
 & \left. + \sum_{\nu=0}^{i-1} \frac{(-1)^{i-\nu} i! (i - \nu)! P(\kappa - m_i, \nu)}{\nu! P(m_i - 2 - i + \nu, i - \nu + 1) (m_i - 1)!} + \frac{P(\kappa - m_i, i)}{(m_i - 1) m_i!} \right) \leq q
 \end{aligned}$$

Enclosure Theorems

Theorem 1. *The recess condition holds for sufficiently large values of κ .*

Theorem 2. *If the recess condition holds, then*

$$|a_k r^k| \leq q \cdot \max_{\nu=0}^{\kappa+n-1} |a_\nu r^\nu| =: A \quad \text{for all } k \geq \kappa + n.$$

Theorem 3. *If the recess condition holds, then for $h = \omega r$, $0 \leq \omega < 1$, and $k \geq \kappa$,*

$$\left| y(h) - \sum_{\nu=0}^{k-1} a_\nu h^\nu \right| = \left| \sum_{\nu=k}^{\infty} a_\nu r^\nu \omega^\nu \right| \leq A \sum_{\nu=k}^{\infty} \omega^\nu = \frac{A \omega^k}{1 - \omega}.$$

Practical calculations

- Step size control:

Recess condition can be tested a priori for r and κ .

- Very large step sizes:

Multiple precision arithmetic required.

Computation of B_i

Analytic function: $p(z) = \sum_{j=0}^{\infty} b_j z^j$, $z \in \mathbb{C}$, $|z| \leq r$.

Cauchy's estimate: $|b_j| \leq \frac{M(r)}{r^j}$, $j \in \mathbb{N}$, $M(r) := \max_{|z|=r} |p(z)|$.

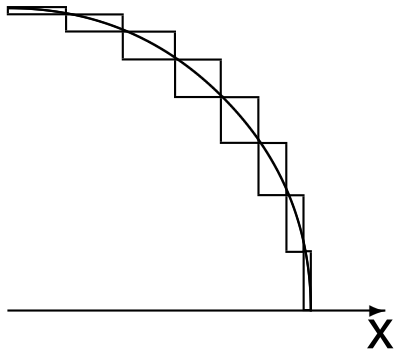
CE of m -th derivative: $p^{(m)}(z) = \sum_{j=0}^m P(j, m) b_{j+m} z^j$

$$\Rightarrow |P(j, m) b_{j+m}| \leq \frac{N(r)}{r^j}, \quad N(r) := \max_{|z|=r} |p^{(m)}(z)|.$$

$$\Leftrightarrow |b_j| \leq \frac{B}{P(j-m, m) r^j}, \quad B = N(r) r^m, \quad \text{for } j \geq m.$$

Interval Computation of M

iy



Cover C with complex intervals

$$Z_k, \quad k = 1, \dots, k_{\max}.$$

Use complex interval functions to compute

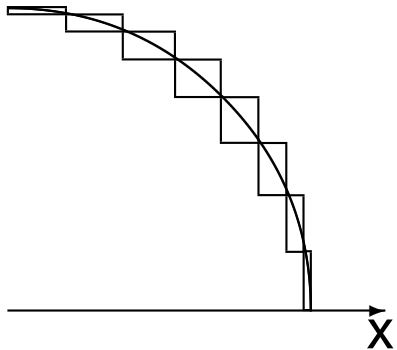
$$[\underline{F}_k, \overline{F}_k] \supseteq |f(Z_k)| \text{ for all } k.$$

$$\Rightarrow M(r) \leq \max_k \overline{F}_k.$$

Adaptive refinement of Z_k with branch and bound algorithm.

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Adaptive refinement of Z_k with branch and bound algorithm.

Example 1: $y'' = y, y(0) = 1, y'(0) = -1.$

Exact Solution: $y(x) = e^{-x}.$

h	$[y](h)$	L	κ	Time
20	2.061 153 622 438 55 ₇ ⁹ E-009	3	97	0.11
40	4.248 354 255 291 5 ₈₈ ⁹¹ E-018	4	170	0.39
100	3.720 075 976 020 83 ₅ ⁷ E-044	7	386	2.47
200	1.383 896 526 736 73 ₇ ⁸ E-087	12	745	12.74
300	5.148 200 222 412 01 ₂ ⁶ E-131	18	1104	37.19

AWA:

h	$[y](h)$	Degree	No. of steps	Time
10	4.539 99 ₂₈ ³² E-05	20	33	0.11
15	3.0 ₅₇ ⁶¹ E-07	20	50	0.17
20	3.136 -2.724 E-08	20	69	0.28
40	1.422 -1.422 E+01	20	101	0.39

Example 2

$$y^{(4)} = (x^2 + 10x + 26) y^{(3)} + (-20x - 99.5) y'' + (x^2 + 10x + 25) y' + (-2x^2 - 4x + 29.5) y$$

$$y(0) = 5, y'(0) = 4, y''(0) = 3, y^{(3)}(0) = 2. \quad (y(x) = (5 - x) e^x.)$$

h	ε_{abs}	$[y](h)$	L	κ	Time
1	—	1.087...61 ₇ ⁹ E+01	2	58	0.11
1.25	—	1.308...19 ₀ ¹ E+01	2	72	0.17
1.5	—	1.568...32 ₂ ³ E+01	2	88	0.22
4	—	5.459...42 ₂ ⁶ E+01	7	364	2.25
5	1E-050	$\begin{matrix} 4.951 \\ -4.951 \end{matrix}$ E-065	12	548	5.28

AWA:

h	$[y](h)$	Degree	No. of steps	Time
1	1.087 312 $\begin{matrix} 77 \\ 68 \end{matrix}$ E+01	20	61	2.42
1.25	1.30 $\begin{matrix} 91 \\ 86 \end{matrix}$ E+01	20	81	3.13
1.5	$\begin{matrix} 34.960 \\ -6.183 \end{matrix}$	20	101	3.95

Example 3

$$y'' = \alpha e^x y + e^{-x} - \alpha, \quad y(0) = 1, \quad y'(0) = -1. \quad (y(x) = e^{-x}.)$$

α	x	$[y](x)$	Prec.	Steps	k_{max}
0.01	5	6.737 ... 46 ₆ ⁹ E-03	2	1	105
0.01	7.5	5.530 843 701 4 ₇₅ ⁸⁴ E-04	2	2	105
0.01	10	6.250 2.846 E-05	2	3	161
10	4	1.831 ... 41 ₇ ⁹ E-02	3	1	201
250	2.25	1.053 ... 64 ₄ ⁴ E-01	3	1	181

AWA:

α	x	$[y](x)$	Order	Steps
0.01	5	6.737 946 999 ₀₆₇ ¹⁰⁹ E-03	20	10
0.01	7.5	5.530 843 ₆₆₁ ⁷⁵³ E-04	20	22
0.01	10	4.336 -3.294 E-03	20	67
10	4	1.070 -0.688 E+02	20	84
250	1.5	1.341 -0.451 E+00	20	78

Nonlinear ODEs

$$y' = f(p(x), y), \quad y(0) = \eta.$$

$$f_i(x) = \sum_{j=1}^J p_{ij}(x) \prod_{l=1}^n y_l^{\alpha_{ijl}}, \quad p_{ij}(x) = \sum_{k=0}^{\infty} b_{ijk} x^k,$$

$$\alpha_{ijl} \in \mathbb{N}_0, \quad b_{ijk} \in \mathbb{R}, \quad |b_{ijk}| \leq \frac{B_{ij}}{R^k}.$$

(f_i is a polynomial in y , with analytic coefficient functions $p_{ij}(x)$).

Power series solution:

$$y_i(x) := \sum_{k=0}^{\infty} a_{ik} x^k, \quad |x| < r \text{ for some } r > 0.$$

Estimation: For $i = 1, \dots, n$ compute $A_i > 0$ and $r > 0$, so that

$$|a_{ik}| \leq \frac{A_i}{r^k} \text{ for all } i = 1, \dots, n \text{ and all } k \in \mathbb{N}.$$

Example 4: Lorenz Equations

$$\dot{x} = -\sigma x + \sigma y, \quad x(0) = 0; \quad \sigma = 10$$

$$\dot{y} = rx - y - xz, \quad y(0) = 1; \quad r = 28$$

$$\dot{z} = -bz + xy, \quad z(0) = 0; \quad b = \frac{8}{3}.$$

Comparison of step sizes:

t	Steps	AWA
0.25	6	12
0.5	13	31
1.0	26	64